

## Abstract

We introduce the class of **priority planning** strategies for suppressing **SIS epidemics** in a network that performs dynamic allocation of treatment resources with limited efficiency to the infected nodes, according to a precomputed **priority-order**. Then, using recent theoretical results that highlight the role of the **maxcut** of a node ordering and

the extinction time of an epidemic, we propose a simple and efficient strategy called **MaxCut Minimization** (MCM) that outperforms competing state-of-the-art strategies in simulated epidemic scenarios that include artificially generated networks as well as real transportation networks.

## BACKGROUND

### 1. Motivation

#### Applications of Diffusion Control in Networks

- *Epidemiology*: Limiting the spread of a disease in a population.
- *Marketing*: Increase product adoption in a community.
- *Rumor spread*: Preventing false rumors from reaching a large audience by providing good information to key users of a social network.

**Related Work**: CURE policy [1], Activity shaping in SNs [2].

### 2. SIS epidemic and control model

#### Notations

Let  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  be an undirected graph of size  $N$ , and  $A$  its adjacency matrix. An **SIS epidemic** is described by the *network state vector*  $X(t) \in \{0, 1\}^N$ . A **control action** is described by a *resource allocation vector*  $\rho(t) \in \mathbb{R}_+^N$ .

#### Diffusion model

Following the formalism of [3], we model an epidemic under a control action as a **stochastic process** with the following transition rates:

$$\begin{aligned} X_i(t) : 0 \rightarrow 1 & \text{ at rate } \beta \sum_j A_{ij} X_j(t); \\ X_i(t) : 1 \rightarrow 0 & \text{ at rate } \delta + \rho_i(t), \end{aligned} \quad (1)$$

where  $\beta \geq 0$ ,  $\delta \geq 0$  are the transmission and recovery rates of the epidemic.

#### Control model

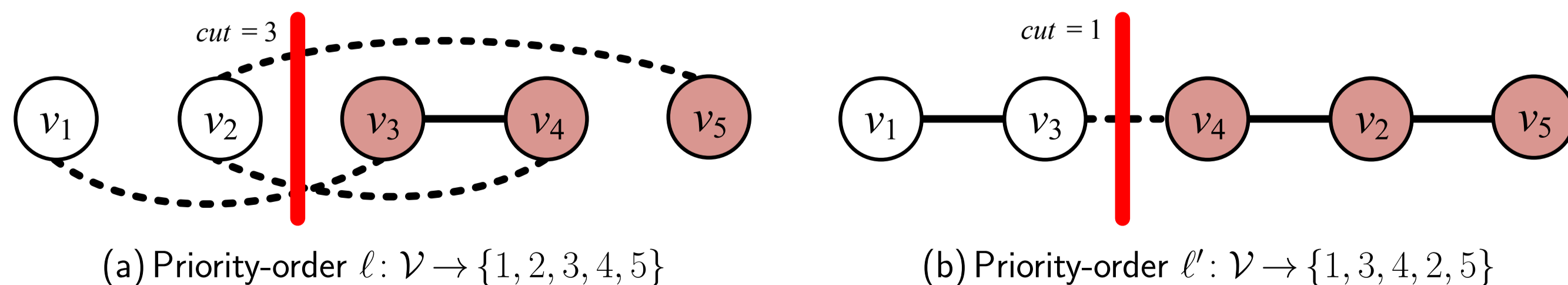
We impose three constraints on the resource allocation vector  $\rho(t)$ :

- *Causality*:  $\rho(t)$  should only depend on past values of  $X(t)$ ,
- *Limited budget*:  $\sum_i \rho_i(t) \leq r$ ,
- *Limited efficiency*:  $\forall i, \rho_i(t) \leq \rho$ .

### 3. Maxcut of a priority-order

A priority-order  $\ell : \mathcal{V} \rightarrow \{1, \dots, N\}$  is a node ordering that describes the order in which an epidemic should be removed from the network, and its maxcut is:

$$C^*(\ell) = \max_{c=1, \dots, N} \sum_{i,j} A_{ij} \mathbb{1}\{\ell(v_i) < c \leq \ell(v_j)\}. \quad (2)$$



**Figure**: Two priority-orders (from left to right) leading to different maxcuts:  $C^*(\ell) = 3$  for (a) and  $C^*(\ell') = 1$  for (b). The cut (vertical red line) separates the nodes in two sets (white are healthy, and red are infected). The second priority-order  $\ell'$  is optimal and the network has a cutwidth  $\mathcal{W} = 1$  (the minimum  $C^*(\ell)$  for any ordering  $\ell$ ).

## THE MCM ALGORITHM

### 4. Priority planning

The introduced approach distributes resources to the **top- $q$  infected nodes** according to a **fixed priority-order**  $\ell$  of the nodes in  $\mathcal{V}$ . The allocated amount of resources should match the available resource budget  $r$ , thus  $q = \min\{\lceil \frac{r}{\rho} \rceil, \sum_i X_i(t)\}$ . The resource allocation vector is then defined as:

$$\rho_i(t) = \begin{cases} \frac{r}{\lceil \frac{r}{\rho} \rceil} & \text{if } X_i(t) = 1 \text{ and } \ell(v_i) \leq \theta(t); \\ 0 & \text{otherwise,} \end{cases} \quad (3)$$

where  $\theta(t)$  is a threshold adjusted s.t.  $\sum_i \mathbb{1}_{\{\rho_i(t) > 0\}} = q$ .

### 5. The MCM algorithm

► *Prior to the diffusion process*: Given a network  $\mathcal{G}$ , we compute, a *priority-order*  $\ell_{MCM}(\mathcal{G})$  with **minimum maxcut**  $C^*(\ell)$ :

$$\ell_{MCM}(\mathcal{G}) = \operatorname{argmin}_{\ell} C^*(\ell). \quad (4)$$

► *During the diffusion process*: MCM distributes the resource budget to the infected nodes according to the order  $\ell_{MCM}(\mathcal{G})$ .

### 6. Practical implementation

#### Relaxation of the optimization problem

Minimizing the maxcut is a very hard combinatorial problem. We relax it to the *Minimum Linear Arrangement* problem which optimizes the meancut:

$$\text{MLA: } \operatorname{argmin}_{\ell} \frac{1}{N} \sum_{c=1, \dots, N} \sum_{i,j} A_{ij} \mathbb{1}\{\ell(v_i) < c \leq \ell(v_j)\}. \quad (5)$$

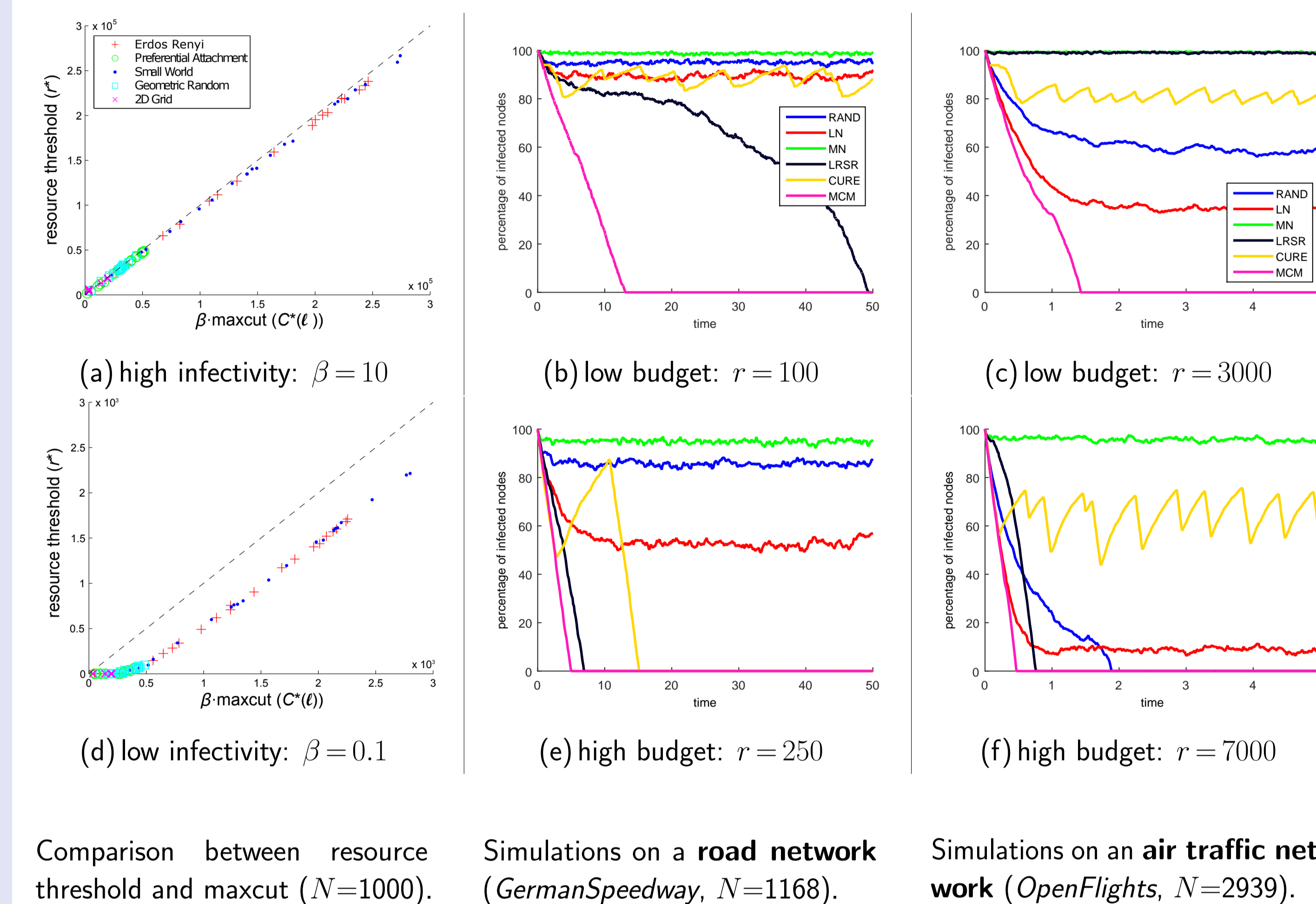
#### Computation of the optimal priority-order

The MLA problem solver that we developed for our simulations follows the steps below and uses a hierarchical approach to take advantage of the group structure of social and contact networks:

- First, we identify dense clusters by applying *spectral clustering* and we order those clusters (considered as high-level nodes) using *spectral sequencing* [4].
- Then, we compute a good ordering of the nodes inside each cluster independently using spectral sequencing followed by an iterative approach which is based on random node swaps (swap heuristics inspired by [5]).
- Finally, we place and orient properly the clusters' node segments and reapply the same swap-based approach as step two to refine the overall ordering.

## EXPERIMENTS

### 7. Control of simulated epidemics



### 8. Discussion on results

- Very good correlation was observed between the maxcut and the resource budget  $r^*$  beyond which the epidemic is suppressed.
- The maxcut can be used as a quality metric for any priority-planning.
- MCM strategy outperforms its competitors in simulated epidemics on real transportation networks.

### References

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